

Algorithm 1 `sample_d1`

Input:
 \mathbf{v} : velocity field
 \bar{t} : anchor time
 $\mathcal{T} = (t_1, \dots, t_L), \mathcal{X} = (x_1, \dots, x_M), \mathcal{Y} = (y_1, \dots, y_N)$: regular grid
 $\mathcal{D}_{1,pos} \leftarrow$ array of 2D vectors, size (L, M, N) \triangleright saves destination of flow map
 $\mathcal{D}_{1,bound} \leftarrow$ array of bytes, size (L, M, N) \triangleright saves boundary flag
 $\mathbf{bounds} \leftarrow ((x_1, x_M), (y_1, y_N))$
 $\mathbf{vc} \leftarrow \text{constrain}(\mathbf{v}, \mathbf{bounds})$ \triangleright sets OutOfBounds outside grid

for $(i, j, k) \in (1, \dots, L) \times (1, \dots, M) \times (1, \dots, N)$ **in parallel do**
 $t \leftarrow \mathcal{T}[i]$
 $\mathbf{x} \leftarrow (\mathcal{X}[j], \mathcal{Y}[k])$
 $\text{state} \leftarrow \text{integrate}(\mathbf{vc}, \mathbf{x}, t, \bar{t})$ \triangleright integrates \mathbf{v} from (\mathbf{x}, t) to \bar{t} , returns end state
 $(\mathbf{p}, b) \leftarrow \text{encode_d1}(\mathbf{bounds}, \text{state})$ \triangleright see Algorithm 4
 $\mathcal{D}_{1,pos}[i, j, k] \leftarrow \mathbf{p}$
 $\mathcal{D}_{1,bound}[i, j, k] \leftarrow b$

end for
return $\mathcal{D}_{1,pos}, \mathcal{D}_{1,bound}$

Algorithm 2 `sample_d2`

Input:
 \mathbf{v} : velocity field
 \bar{t} : anchor time
 $\mathcal{T} = (t_1, \dots, t_L), \mathcal{X} = (x_1, \dots, x_M), \mathcal{Y} = (y_1, \dots, y_N)$: regular grid
 $\mathcal{D}_2 \leftarrow$ array of 2D vectors, size (L, M, N)

for $(i, j, k) \in (1, \dots, L) \times (1, \dots, M) \times (1, \dots, N)$ **in parallel do**
 $t \leftarrow \mathcal{T}[i]$
 $\mathbf{x} \leftarrow (\mathcal{X}[j], \mathcal{Y}[k])$
 $\text{state} \leftarrow \text{integrate}(\mathbf{v}, \mathbf{x}, \bar{t}, t)$ \triangleright integrates from (\mathbf{x}, t) to \bar{t} , returns end state
 $\mathcal{D}_2[i, j, k] \leftarrow \text{pos}(\text{state})$ \triangleright may be undefined vector if leaving domain

end for
return \mathcal{D}_2

Algorithm 3 `sample_d3`

Input:
 \mathbf{v} : velocity field
 \bar{t} : anchor time
 $\mathcal{T} = (t_1, \dots, t_L), \mathcal{X} = (x_1, \dots, x_M), \mathcal{Y} = (y_1, \dots, y_N)$: regular grid
 b : boundary flag
 $\text{fixed} \leftarrow (x_1, y_1, x_M, y_N)[b]$
 $\text{dim} \leftarrow b \bmod 2$
 $K \leftarrow \text{dim} = 1 ? N : M$ \triangleright length of free dimension
 $\mathcal{D}_{3,b} \leftarrow$ array of 2D vectors, size (L, L, K) \triangleright grid with two time and one space dimensions

for $(i, j, k) \in (1, \dots, L) \times (1, \dots, L) \times (1, \dots, K)$ **in parallel do**
 $\mathbf{x} \leftarrow \text{dim} = 1 ? (\text{fixed}, \mathcal{Y}[k]) : (\mathcal{X}[k], \text{fixed})$ \triangleright construct position on the boundary
 $\text{state} \leftarrow \text{integrate}(\mathbf{v}, \mathbf{x}, \mathcal{T}[i], \mathcal{T}[j])$
 $\mathcal{D}_{3,b}[i, j, k] \leftarrow \text{pos}(\text{state})$

end for
return $\mathcal{D}_{3,b}$

1. Algorithms

Here we provide some information about the algorithms. Please see Sections 4.2 (Sampling and Reconstruction) and 4.4 (Implementation) of the main paper first. Again, we describe the 2D case for simplicity, but the 3D extension is straightforward.

Algorithms 1, 2 and 3 show pseudocode for sampling the fields $\mathbf{d}_1, \mathbf{d}_2$ and \mathbf{d}_3 . Algorithm 4 explains how we encode the result of \mathbf{d}_1 . If the pathline stays in the domain, the result is $(\mathbf{x}', 0)$, i.e., the spatial end position and a zero. Latter encodes that we will use \mathbf{d}_2 in the reconstruction. Otherwise, the integration ended at a spatial boundary and we search the respective side (x or y , min or max). Each boundary side has a flag $b \in \{1, 2, 3, 4\}$. Additionally, we return the vector \mathbf{p} which contains the free values, i.e., the end time and the remaining spatial value.

When computing \mathbf{d}_3 we only need to call `sample_d3` for a given boundary flag if $b \in \mathcal{D}_{1,bound}$, i.e., at least one pathline left the domain on this side. This saves time and memory usage.

Algorithm 4 `encode_d1`

Input:
 $\mathbf{bounds} = ((x_s, x_e), (y_s, y_e))$: domain boundary
 state : contains information of the end state of an integration
 $t' \leftarrow \text{time}(\text{state})$
 $\mathbf{x}' \leftarrow \text{pos}(\text{state})$
 $\mathbf{v}' \leftarrow \text{velocity}(\text{state})$ \triangleright important: vector is inverted if applying backward integration

if $\text{succeeded}(\text{state})$ **then** \triangleright default case: pathline stayed inside domain
return $(\mathbf{x}', 0)$

end if
 $b \leftarrow -1$ \triangleright boundary flag
 $d_{min} \leftarrow \infty$ \triangleright current min distance to boundary

for $i \in \{1, 2\}$ **do** \triangleright search for boundary flag b where the domain was left
 $d_s \leftarrow \mathbf{x}'[i] - \mathbf{bounds}[i][1]$ \triangleright distance to x_s or y_s
 $d_e \leftarrow \mathbf{bounds}[i][2] - \mathbf{x}'[i]$ \triangleright distance to x_e or y_e
if $(\mathbf{v}'[i] < 0) \wedge (d_s < d_{min})$ **then**
 $d_{min} \leftarrow d_s, b \leftarrow i$
else if $(\mathbf{v}'[i] > 0) \wedge (d_e < d_{min})$ **then**
 $d_{min} \leftarrow d_e, b \leftarrow 2 + i$
end if

end for
 $\text{dim} \leftarrow b \bmod 2$
 $\mathbf{p} \leftarrow (\text{dim} = 1) ? (t', y') : (t', x')$ \triangleright compose time-space vector

return (\mathbf{p}, b)

Algorithm 5 `reconstruct`

Input:
 afm : anchor flow map with grid and pre-sampled data
 \mathbf{x}_0, t_0, t_1 : evaluation input
 $(\mathcal{T}, \mathcal{X}, \mathcal{Y}) \leftarrow \text{grid_data}(\text{afm})$
 $\mathcal{D}_{1,pos}, \mathcal{D}_{1,bound} \leftarrow \text{get_d1}(\text{afm})$
 $\mathbf{x}_1 \leftarrow (0, 0)$ \triangleright this will be the result
 $(\mathcal{I}, \mathcal{W}) \leftarrow \text{interpolants}((\mathcal{T}, \mathcal{X}, \mathcal{Y}), (t_0, \mathbf{x}_0))$ \triangleright get indices and weights for interpolation

for $(i, w) \in (\mathcal{I}, \mathcal{W})$ **do**
 $(\mathbf{p}, b) \leftarrow (\mathcal{D}_{1,pos}[i], \mathcal{D}_{1,bound}[i])$
if $b = 0$ **then**
 $\mathcal{D}_2 \leftarrow \text{get_d2}(\text{afm})$
 $\mathbf{x}_1 \leftarrow \mathbf{x}_1 + w \cdot \text{interpolate}((\mathcal{T}, \mathcal{X}, \mathcal{Y}), \mathcal{D}_2, (t_1, \mathbf{p}))$
else
 $\mathcal{D}_{3,b} \leftarrow \text{get_d3}(\text{afm}, b)$
 $(t', s') = \mathbf{p}$ \triangleright free time and space values
 $\mathbf{q} \leftarrow (t', t_1, s')$ \triangleright insert target time t_1
 $\text{dim} \leftarrow b \bmod 2$
 $\text{grid} \leftarrow \text{dim} = 1 ? (\mathcal{T}, \mathcal{T}, \mathcal{Y}) : (\mathcal{T}, \mathcal{T}, \mathcal{X})$ \triangleright determine corresponding grid
 $\mathbf{x}_1 \leftarrow \mathbf{x}_1 + w \cdot \text{interpolate}(\text{grid}, \mathcal{D}_{3,b}, \mathbf{q})$
end if

end for
return \mathbf{x}_1

For Algorithm 5, note the difference between functions `interpolants` and `interpolate`. First returns vertex indices and weights for multilinear interpolations in the respective grid, latter directly applies interpolation with provided data \mathcal{D}_2 or $\mathcal{D}_{3,b}$.

2. Error Plots

The main paper shows error plots of the Double Gyre dataset. The ABC flow (Figure 2) provides similar plots for mean errors. In case of maximum errors, however, the median and 25% quartile results of AFM_c are worse than for FFM, even with res_A . The reason: all six boundary surfaces must be sampled for \mathbf{d}_3 . This decreases the factor between res_F and res_A , leading to a smaller advantage.

The error plots for the Cylinder and Tangaroa datasets (Figures 1 and 3) look much less „regular”. This stems from pathlines leaving the domain at some points which have a strong impact on the averages. This is especially prominent since the pathlines which stay in the domain longest show strong turbulence, and such pathlines are the hardest ones to approximate. For globally defined datasets like the Double Gyre and the ABC flow this does not apply.

Dataset	Setup	Method	Mean errors	Max errors	FP	FN
Double Gyre	res_F	FFM	0.012 ± 0.017	0.58 ± 0.59	0%	0%
		AFM _c	0.014 ± 0.017	0.74 ± 0.52	0%	0%
	res_A	AFM _c	0.0028 ± 0.0042	0.54 ± 0.51	0%	0%
Cylinder	res_F	FFM	0.090 ± 0.14	1.48 ± 1.50	0.0064%	1.308%
		AFM _c	0.093 ± 0.14	1.59 ± 1.49	0.0029%	1.764%
	res_A	AFM _s	0.042 ± 0.046	3.18 ± 1.90	0.0034%	1.090%
		AFM _e	0.039 ± 0.057	1.96 ± 1.85	0.0031%	0.684%
ABC	res_F	FFM	0.21 ± 0.25	3.70 ± 4.20	0%	0%
		AFM _c	0.25 ± 0.28	4.45 ± 3.90	0%	0%
	res_A	AFM _c	0.13 ± 0.16	3.99 ± 3.70	0%	0%
Tangaroo	res_F	FFM	0.027 ± 0.046	0.29 ± 0.21	0.0035%	2.640%
		AFM _c	0.024 ± 0.042	0.30 ± 0.19	0.0017%	3.580%
	res_A	AFM _c	0.022 ± 0.038	0.31 ± 0.19	0.0014%	2.447%

Table 1: Error statistics of the AFM and FFM compared to the ground truth (i.e., pathline integration). For the AFM method, the subscript determines which anchor value \tilde{t} is selected, with values t_s for AFM_s, t_e for AFM_e, and the center value for AFM_c. As explained in the main paper, we tracked the mean and maximum errors for each (t_0, t_1) sample. The values here show the respective means plus/minus standard deviations. Further, FP and NP note percentages of false positives and false negatives, i.e., falsely returning a flow map value or falsely returning none.

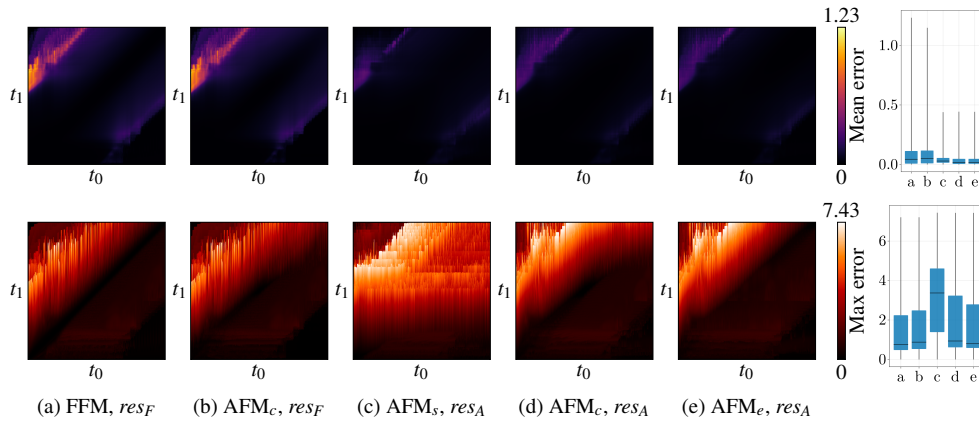


Figure 1: Error plots for the Cylinder dataset.

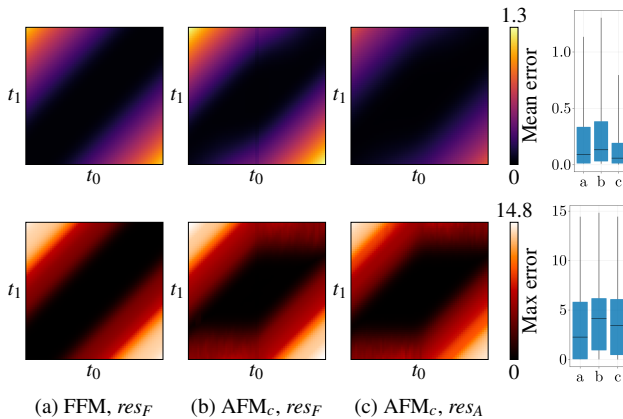


Figure 2: Error plots for the ABC flow dataset.

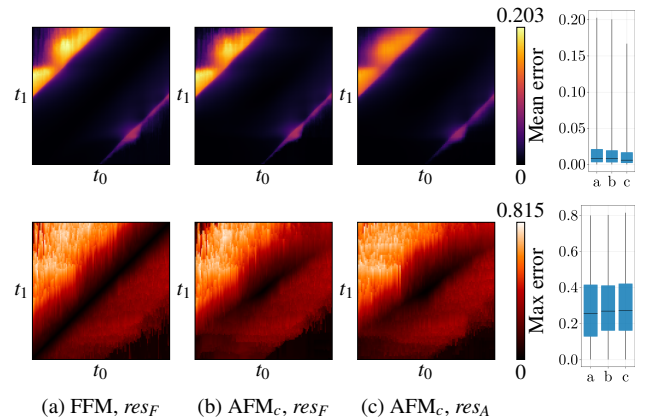


Figure 3: Error plots for the Tangaroo dataset.